

③  
MRC Technical Summary Report #2732

THE LINDSTROM-MADDEN METHOD FOR SERIES  
SYSTEMS WITH REPEATED COMPONENTS

Andrew P. Soms

**AD-A147 238**

**Mathematics Research Center  
University of Wisconsin—Madison  
610 Walnut Street  
Madison, Wisconsin 53705**

August 1984

(Received July 25, 1984)

DTIC  
JUL 25 1984  
D

**Approved for public release  
Distribution unlimited**

Sponsored by

U. S. Army Research Office  
P. O. Box 12211  
Research Triangle Park  
North Carolina 27709

Office of Naval Research  
800 North Quincy Street  
Arlington, VA 22217

84 11 06

041  
221

UNIVERSITY OF WISCONSIN-MADISON  
MATHEMATICS RESEARCH CENTER

THE LINDSTROM-MADDEN METHOD FOR SERIES SYSTEMS WITH REPEATED COMPONENTS

Andrew P. Soms \*

Technical Summary Report #2732  
August 1984

ABSTRACT

The Lindstrom-Madden method of computing lower confidence limits for series systems with unlike components is extended to series systems with repeated components utilizing the results of Harris and Soms (1983). An exact solution is given for no failures and key test results, together with an approximation for the general case. Numerical examples are also provided.

DMS (MOS) Subject Classifications: 62N05, 90B25

Key Words: Lindstrom-Madden approximation; Optimal confidence limits;  
Reliability; Repeated components; Series system

Work Unit Number 4 (Statistics and Probability)

---

\*Department of Mathematical Sciences, University of Wisconsin-Milwaukee, Milwaukee, WI 53201.

---

Sponsored in part by the United States Army under Contract No. DAAG29-80-C-0041, the Office of Naval Research under Contract No. N00014-79-C-0321, and the University of Wisconsin-Milwaukee.

# SIGNIFICANCE AND EXPLANATION

Series systems with repeated components arise often in engineering and physics. It is therefore important to utilize data obtained on individual components in an efficient manner when assessing the reliability of the combined system. This paper gives one method for doing so.

Application For	
1. Name of Component	<input checked="" type="checkbox"/>
2. Name of Component	<input type="checkbox"/>
3. Name of Component	<input type="checkbox"/>
4. Name of Component	<input type="checkbox"/>
5. Name of Component	<input type="checkbox"/>
6. Name of Component	<input type="checkbox"/>
7. Name of Component	<input type="checkbox"/>
8. Name of Component	<input type="checkbox"/>
9. Name of Component	<input type="checkbox"/>
10. Name of Component	<input type="checkbox"/>
11. Name of Component	<input type="checkbox"/>
12. Name of Component	<input type="checkbox"/>
13. Name of Component	<input type="checkbox"/>
14. Name of Component	<input type="checkbox"/>
15. Name of Component	<input type="checkbox"/>
16. Name of Component	<input type="checkbox"/>
17. Name of Component	<input type="checkbox"/>
18. Name of Component	<input type="checkbox"/>
19. Name of Component	<input type="checkbox"/>
20. Name of Component	<input type="checkbox"/>
21. Name of Component	<input type="checkbox"/>
22. Name of Component	<input type="checkbox"/>
23. Name of Component	<input type="checkbox"/>
24. Name of Component	<input type="checkbox"/>
25. Name of Component	<input type="checkbox"/>
26. Name of Component	<input type="checkbox"/>
27. Name of Component	<input type="checkbox"/>
28. Name of Component	<input type="checkbox"/>
29. Name of Component	<input type="checkbox"/>
30. Name of Component	<input type="checkbox"/>
31. Name of Component	<input type="checkbox"/>
32. Name of Component	<input type="checkbox"/>
33. Name of Component	<input type="checkbox"/>
34. Name of Component	<input type="checkbox"/>
35. Name of Component	<input type="checkbox"/>
36. Name of Component	<input type="checkbox"/>
37. Name of Component	<input type="checkbox"/>
38. Name of Component	<input type="checkbox"/>
39. Name of Component	<input type="checkbox"/>
40. Name of Component	<input type="checkbox"/>
41. Name of Component	<input type="checkbox"/>
42. Name of Component	<input type="checkbox"/>
43. Name of Component	<input type="checkbox"/>
44. Name of Component	<input type="checkbox"/>
45. Name of Component	<input type="checkbox"/>
46. Name of Component	<input type="checkbox"/>
47. Name of Component	<input type="checkbox"/>
48. Name of Component	<input type="checkbox"/>
49. Name of Component	<input type="checkbox"/>
50. Name of Component	<input type="checkbox"/>
51. Name of Component	<input type="checkbox"/>
52. Name of Component	<input type="checkbox"/>
53. Name of Component	<input type="checkbox"/>
54. Name of Component	<input type="checkbox"/>
55. Name of Component	<input type="checkbox"/>
56. Name of Component	<input type="checkbox"/>
57. Name of Component	<input type="checkbox"/>
58. Name of Component	<input type="checkbox"/>
59. Name of Component	<input type="checkbox"/>
60. Name of Component	<input type="checkbox"/>
61. Name of Component	<input type="checkbox"/>
62. Name of Component	<input type="checkbox"/>
63. Name of Component	<input type="checkbox"/>
64. Name of Component	<input type="checkbox"/>
65. Name of Component	<input type="checkbox"/>
66. Name of Component	<input type="checkbox"/>
67. Name of Component	<input type="checkbox"/>
68. Name of Component	<input type="checkbox"/>
69. Name of Component	<input type="checkbox"/>
70. Name of Component	<input type="checkbox"/>
71. Name of Component	<input type="checkbox"/>
72. Name of Component	<input type="checkbox"/>
73. Name of Component	<input type="checkbox"/>
74. Name of Component	<input type="checkbox"/>
75. Name of Component	<input type="checkbox"/>
76. Name of Component	<input type="checkbox"/>
77. Name of Component	<input type="checkbox"/>
78. Name of Component	<input type="checkbox"/>
79. Name of Component	<input type="checkbox"/>
80. Name of Component	<input type="checkbox"/>
81. Name of Component	<input type="checkbox"/>
82. Name of Component	<input type="checkbox"/>
83. Name of Component	<input type="checkbox"/>
84. Name of Component	<input type="checkbox"/>
85. Name of Component	<input type="checkbox"/>
86. Name of Component	<input type="checkbox"/>
87. Name of Component	<input type="checkbox"/>
88. Name of Component	<input type="checkbox"/>
89. Name of Component	<input type="checkbox"/>
90. Name of Component	<input type="checkbox"/>
91. Name of Component	<input type="checkbox"/>
92. Name of Component	<input type="checkbox"/>
93. Name of Component	<input type="checkbox"/>
94. Name of Component	<input type="checkbox"/>
95. Name of Component	<input type="checkbox"/>
96. Name of Component	<input type="checkbox"/>
97. Name of Component	<input type="checkbox"/>
98. Name of Component	<input type="checkbox"/>
99. Name of Component	<input type="checkbox"/>
100. Name of Component	<input type="checkbox"/>

A-1



The responsibility for the wording and views expressed in this descriptive summary lies with MRC, and not with the author of this report.

# THE LINDSTROM-MADDEN METHOD FOR SERIES SYSTEMS WITH REPEATED COMPONENTS

Andrew P. Soms\*

## 1. INTRODUCTION AND SUMMARY

A problem of substantial importance to practitioners in reliability is the statistical estimation of the reliability of a series system of stochastically independent components when some components are repeated, using experimental data collected on the individual components. In the situations discussed in this paper, the component data consist of a sequence of Bernoulli trials. Thus, for component  $i$ ,  $i = 1, 2, \dots, k$ , the data is the pair  $(n_i, Y_i)$ , where  $n_i$  is the number of trials and  $Y_i$  is the number of observations for which the component functions.  $Y_1, Y_2, \dots, Y_k$  are assumed to be mutually independent random variables. We assume that there are  $Y_i$  components of type  $i$ ,  $1 \leq i \leq k$ . Then the parameter of interest is  $h(p_1, p_2, \dots, p_k) = h(\tilde{p})$ , the reliability of the system, where

$$h(\tilde{p}) = \prod_{i=1}^k p_i^{Y_i}.$$

More specifically, it is desired to obtain a Buehler (1957) optimal lower  $1 - \alpha$  confidence limit on  $h(\tilde{p})$ .

The case of  $Y_1 = Y_2 = \dots = Y_k = 1$  has been treated in Sudakov (1974), Winterbottom (1974), and Harris and Soms (1983).

In Section 2 we summarize the general theory of Harris and Soms (1983) applicable here. In Section 3 the exact solutions to no failures and key test results are given. Lindstrom-Madden type approximations are given in Section 4. Section 5 contains numerical examples.

---

\*Department of Mathematical Sciences, University of Wisconsin-Milwaukee, Milwaukee, WI 53201.

---

Sponsored in part by the United States Army under Contract No. DAAG29-80-C-0041, the Office of Naval Research under Contract No. N00014-79-C-0321, and the University of Wisconsin-Milwaukee.

## 2. BUEHLER'S METHOD FOR OPTIMAL CONFIDENCE LIMITS

We now specialize the general results of Harris and Soms (1983) on optimal confidence limits for system reliability to a series system with independent and repeated components. As in Section 1, let

$$h(\tilde{p}) = \prod_{i=1}^k p_i^{y_i},$$

$0 < p_i < 1$ ,  $X_i = n_i - Y_i$ ,  $x_i = n_i - y_i$ ,  $1 \leq i \leq k$ ,  $S = \{\tilde{x} | x_i = 0, 1, \dots, n_i, 1 \leq i \leq k\}$  and let  $g(\tilde{x}) = (x_1, x_2, \dots, x_k)$  be an ordering function, i.e., for real  $x_i$ ,  $0 \leq x_i \leq n_i$ ,  $g(\tilde{x})$  is non-decreasing in each component. It is often convenient to normalize  $g(\tilde{x})$  by letting  $g(\tilde{0}) = 1$  and  $g(\tilde{n}) = 0$ . With such a normalization,  $g(\tilde{x})$  is often selected to be a point estimator of  $h(\tilde{p})$ . Also let  $R = \{r_1, r_2, \dots, r_s, s \geq 2\}$  be the range set of  $g(\tilde{x})$ . With no loss of generality we order  $R$  so that  $r_1 > r_2 > \dots > r_s$  and let  $A_i = \{\tilde{x} | g(\tilde{x}) = r_i, \tilde{x} \in S, i = 1, 2, \dots, s\}$ . The sets  $A_i$  constitute a partition of  $S$  induced by  $g(\tilde{x})$ . We assume throughout that the data is distributed by

$$\begin{aligned} f(\tilde{x}; \tilde{p}) &= P_{\tilde{p}}(\tilde{X} = \tilde{x}) = \prod_{i=1}^k \binom{n_i}{x_i} p_i^{x_i} q_i^{n_i - x_i} \\ &= \prod_{i=1}^k \binom{n_i}{y_i} p_i^{y_i} q_i^{n_i - y_i}, \end{aligned} \quad (2.1)$$

where  $q_i = 1 - p_i$ ,  $i = 1, 2, \dots, k$ . With no loss of generality, we assume

$$n_1 \leq n_2 \leq \dots \leq n_k.$$

From these definitions, it follows that

$$P_{\tilde{p}}\left\{X \in \bigcup_{i=1}^j A_i\right\} = P_{\tilde{p}}\{g(\tilde{X}) \geq r_j\}. \quad (2.2)$$

From (2.1) and (2.2), we have

$$P_{\tilde{p}}\{g(\tilde{X}) > r_j\} = \sum_{i_1=0}^{u_1} \sum_{i_2=0}^{u_2} \cdots \sum_{i_k=0}^{u_k} f(\tilde{i}; \tilde{p}), \quad (2.3)$$

where  $\tilde{i} = (i_1, i_2, \dots, i_k)$  and  $u_2 = u_2(i_1), \dots, u_k = u_k(i_1, i_2, \dots, i_{k-1})$  are integers determined by  $r_j$ . Equivalently,

$$P_{\tilde{p}}\{g(\tilde{X}) > r_j\} = \sum_{i_1=0}^{[t_1]} \sum_{i_2=0}^{[t_2]} \cdots \sum_{i_k=0}^{[t_k]} f(\tilde{i}; \tilde{p}), \quad (2.4)$$

where  $t_2 = t_2(i_1), \dots, t_k = t_k(i_1, i_2, \dots, i_{k-1})$ , with  $t_1 = \sup\{t | 0 \leq t \leq n_1 \text{ and } g(t, 0, 0, \dots, 0) > r_j\}$  and  $t_\ell(i_1, i_2, \dots, i_{\ell-1}) = \sup\{t | 0 \leq t \leq n_\ell \text{ and } g(i_1, i_2, \dots, i_{\ell-1}, t, 0, \dots, 0) > r_j\}$ ,  $\ell = 2, 3, \dots, k$ .

We now introduce the notion of Buehler optimal confidence limits. Let  $g(x) = r_j$ . Then define

$$a_{g(\tilde{x})} = \inf\{h(\tilde{p}) | P_{\tilde{p}}\{\tilde{i} | g(\tilde{i}) > g(\tilde{x})\} > \alpha\}. \quad (2.5)$$

Equivalently, by (2.2), we can also write

$$a_{g(\tilde{x})} = \inf\{h(\tilde{p}) | P_{\tilde{p}}\{X \in \bigcup_{i=1}^j A_i\} > \alpha\}. \quad (2.6)$$

Then we have, from Harris and Soms (1983),

**Theorem 2.1.**  $a_{g(\tilde{x})}$  is a  $1 - \alpha$  lower confidence limit for  $h(\tilde{p})$ . If  $b_{g(\tilde{X})}$  is any other  $1 - \alpha$  lower confidence limit for  $h(\tilde{p})$  with  $b_{r_1} > b_{r_2} > \dots > b_{r_j}$ , then  $b_{g(\tilde{x})} \leq a_{g(\tilde{x})}$  for all  $\tilde{x} \in S$ .

Two possible choices of  $g(\tilde{x})$  are

$$g(\tilde{x}) = \prod_{i=1}^k ((n_i - x_i)/n_i)^{Y_i}, \quad (2.7)$$

or

$$g(\tilde{x}) = \prod_{i=1}^k \prod_{j=0}^{Y_i-1} \left( \frac{n_i - x_i - j}{n_i - j} \right). \quad (2.8)$$

Both reduce to the generally used  $g(\tilde{x})$  for series systems with independent components when  $Y_1 = Y_2 = \dots = Y_k = 1$ , i.e.,

$$g(\tilde{x}) = \prod_{i=1}^k (n_i - x_i) / n_i.$$

Since (2.7) is the maximum likelihood estimator of  $h(\tilde{p})$  we will use it here and from now on it will be understood that  $g(\tilde{x})$  is given by (2.7). With this choice of  $g(\tilde{x})$ , we assume from now on that  $0 < x_i < n_i$ ,  $i = 1, 2, \dots, k$ , since  $a_{g(\tilde{x})} = 0$  if some  $x_i = n_i$ . With this assumption, the  $t_i$  in (2.4) are given by

$$t_1 = n_1 - \left( \prod_{i=1}^k (n_i - x_i)^{Y_i} / \prod_{i=2}^k n_i^{Y_i} \right)^{1/Y_1} \quad (2.9)$$

and

$$t_\ell = n_1 - \left( \prod_{i=1}^k (n_i - x_i)^{Y_i} / \prod_{s=1}^{\ell-1} (n_s - i_s)^{Y_s} \prod_{i=\ell+1}^k n_i^{Y_i} \right)^{1/Y_\ell}, \quad (2.10)$$

$\ell = 2, \dots, k$ , with  $\prod_{i=k+1}^k n_i^{Y_i} = 1$ .

For the purpose of simplifying the calculation of  $a_{g(\tilde{x})}$  in special cases it is necessary to state additional results from Harris and Soms (1983).

**Theorem 2.2.** Let  $g(\tilde{x}) = r_j$  and let

$$f^*(x; a) = \sup_{h(p)=a} P_P\{g(\tilde{X}) > r_j\}, \quad 0 < a < 1. \quad (2.11)$$

Then

$$\inf_{0 < a < 1} f^*(\tilde{x}; a) = 0, \quad \sup_{0 < a < 1} f^*(\tilde{x}; a) = 1$$

and  $f^*(\tilde{x}; a)$  is strictly increasing in  $a$ .

Theorem 2.3.  $f^*(\tilde{x}; a) = \alpha$  has exactly one solution  $a_\alpha$  in  $a$  and  $a_\alpha = a_g(\tilde{x})$ .

### 3. EXACT SOLUTIONS FOR ZERO FAILURES AND KEY TEST RESULTS

We first assume that  $\tilde{x} = (0, 0, \dots, 0) = \tilde{0}$  and use Theorem 2.3 to obtain  $a_g(\tilde{0})$ .

Theorem 3.1. If  $\tilde{x} = \tilde{0}$ , then

$$f^*(\tilde{0}; a) = \sup_{\prod_{i=1}^k p_i = a} \prod_{i=1}^k \frac{n_i}{p_i} = a^{n_j/\gamma_j}, \quad (3.1)$$

where  $n_j/\gamma_j = \min_{1 \leq i \leq k} n_i/\gamma_i$  and

$$a_g(\tilde{0}) = a^{n_j/\gamma_j}. \quad (3.2)$$

Proof.

$$\begin{aligned} \prod_{i=1}^k \frac{n_i}{p_i} &= \left( \prod_{i=1}^k \frac{\gamma_i}{p_i} \right)^{n_j/\gamma_j} \prod_{i=1, i \neq j}^k \frac{(n_i \gamma_j - n_j \gamma_i)/\gamma_j}{p_i} \\ &\leq a^{n_j/\gamma_j}, \end{aligned}$$

since  $n_i \gamma_j - n_j \gamma_i > 0$  is equivalent to  $n_i/\gamma_i > n_j/\gamma_j$ , which is true, and therefore

$\prod_{i=1, i \neq j}^k \frac{(n_i \gamma_j - n_j \gamma_i)/\gamma_j}{p_i} < 1$ . (3.1) follows by noting that the choice  $p_j = a^{1/\gamma_j}$ ,  $p_i = 1$ ,

$i \neq j$ , gives  $\prod_{i=1}^k \frac{n_i}{p_i} = a^{n_j/\gamma_j}$ . Then, using Theorem 2.3, we obtain (3.2), which reduces to the known series result if  $\gamma_1 = \gamma_2 = \dots = \gamma_k = 1$ .



We now turn to analogues of key test results (see, e.g., Winterbottom (1974) and Harris and Soms (1983)). We define a key test result if  $\gamma_1 = \max_{1 \leq i \leq k} \gamma_i$  (recall that  $n_1 = \min_{1 \leq i \leq k} n_i$ ) and  $\tilde{x} = (x_1, 0, \dots, 0)$ .

Theorem 3.2. If  $\tilde{x}$  is a key test result and

$$\left\{ \tilde{z} \mid \prod_{i=1}^k (n_i - z_i)^{\gamma_i} > \prod_{i=1}^k (n_i - x_i)^{\gamma_i} \right\} = \left\{ \tilde{z} \mid \sum_{i=1}^k (n_i - z_i) > \sum_{i=1}^k (n_i - x_i) \right\}, \quad (3.3)$$

then

$$f^*(\tilde{x}; a) = I_{1/\gamma_1}(n - x_1, x_1 + 1), \quad (3.4)$$

where  $I_x(a, b)$  is the incomplete beta function. Let  $b_\alpha$  denote the solution in  $b$  of

$$\alpha = I_b(n_1 - x_1, x_1 + 1).$$

Then  $a_{g(\tilde{x})} = b_\alpha^{\gamma_1}$ . Note that  $b_\alpha$  is the usual  $1 - \alpha$  lower confidence limit on  $p$ , given  $x_1$  failures in  $n_1$  trials.

Proof. Without loss of generality we can assume that  $n_1 = n_2 = \dots = n_k$ , for otherwise we can write (2.4) as

$$\begin{aligned} P_p\{g(\tilde{X}) > r_j\} &= \sum_{i_1=0}^{x_1} \binom{n_1}{i_1} p_1^{n_1-i_1} q_1^{i_1} \sum_{i_2=0}^{x_1-i_1} \binom{n_2}{i_2} p_2^{n_2-i_2} q_2^{i_2} \dots \\ &\quad \sum_{i_{k-1}=0}^{x_1-i_1-i_2-\dots-i_{k-2}} \binom{n_{k-1}}{i_{k-1}} p_{k-1}^{n_{k-1}-i_{k-1}} q_{k-1}^{i_{k-1}} I_{p_k}(n_k - \\ &\quad (x_1-i_1-i_2-\dots-i_{k-1}), x_1-i_1-i_2-\dots-i_{k-1}+1) \\ &\leq \sum_{i_1=0}^{x_1} \binom{n_1}{i_1} p_1^{n_1-i_1} q_1^{i_1} \dots \sum_{i_{k-1}=0}^{x_1-i_1-i_2-\dots-i_{k-2}} \binom{n_{k-1}}{i_{k-1}} p_{k-1}^{n_{k-1}-i_{k-1}} q_{k-1}^{i_{k-1}} I_{p_k}(n_1 - \\ &\quad (x_1-i_1-i_2-\dots-i_{k-1}), x_1-i_1-i_2-\dots-i_{k-1}+1), \end{aligned} \quad (3.5)$$

where  $g(\tilde{x}) = r_j$ , by the monotone likelihood ratio property of the beta distribution  $(I_x(a,b))$  has a monotone likelihood ratio in  $-a$  for fixed  $b$ , which implies that  $I_x(a,b)$  is a decreasing function of  $a$ . A similar argument applies to the other indexes. Thus, if (3.4) is true for  $n_1 = n_2 = \dots = n_k$ , by (3.5) it follows for  $n_1 \leq n_2 \leq \dots \leq n_k$ .

So, assuming  $\tilde{n} = (n_1, n_1, \dots, n_1)$ , we seek to maximize

$$P_p \left\{ \bigwedge_{i=1}^k \bigwedge_{j=1}^{n_1} Y_{ij} \geq \bigwedge_{i=1}^k (n_i - x_i) = \bigwedge_{i=1}^k y_i \right\}, \quad (3.6)$$

where  $Y_{ij}$  are independent Bernoulli random variables with parameter  $p_i$  and  $\prod_{i=1}^k p_i^{y_i} = a$ . If  $\prod_{i=1}^k p_i^{y_i} = a$ , then  $\prod_{i=1}^k p_i$  ranges from  $a^{1/y_j}$  to

$a^{1/y_1}$ ,  $y_j = \min_{1 \leq i \leq k} y_i$ . This is seen as follows:

$$\begin{aligned} \prod_{i=1}^k p_i &= \left( \prod_{i=1}^k p_i^{y_i} \right)^{1/y_1} \prod_{i=2}^k p_i^{1-y_i/y_1} \\ &= a^{1/y_1} \prod_{i=2}^k p_i^{(y_1-y_i)/y_1} \leq a^{1/y_1} \end{aligned}$$

and

$$\begin{aligned} \prod_{i=1}^k p_i &= \left( \prod_{i=1}^k p_i^{y_i} \right)^{1/y_j} \prod_{i \neq j}^k p_i^{1-y_i/y_j} \\ &= a^{1/y_j} \prod_{i \neq j}^k p_i^{(y_j-y_i)/y_j} \geq a^{1/y_j} \end{aligned}$$

and the choices  $p_1 = a^{1/y_1}$ ,  $p_2 = \dots = p_k = 1$ , and  $p_j = a^{1/y_j}$ ,  $p_i = 1$ ,  $i \neq j$ , attain these values. From the results of Pledger and Proschan (1971), for each  $b = \prod_{i=1}^k p_i$ ,

$a^{1/\gamma_j} < b < a^{1/\gamma_1}$ , (3.6) is maximized by  $p_1 = b$ ,  $p_i = 1$ ,  $2 \leq i \leq k$ . Further, the maximum over  $b$ ,  $a^{1/\gamma_j} < b < a^{1/\gamma_1}$ , of the maxima for each  $b$  is given by  $p_1 = a^{1/\gamma_1}$ ,  $p_i = 1$ ,  $2 \leq i \leq k$ , by the monotone likelihood ratio property of the binomial distribution, and  $p_1 = a^{1/\gamma_1}$ ,  $p_i = 1$ ,  $2 \leq i \leq k$ , satisfies  $\prod_{i=1}^k p_i^{\gamma_i} = a$ . This completes the proof.

If  $\gamma_1 = \gamma_2 = \dots = \gamma_k = 1$ , some guidelines for the verification of (3.3) are given in Harris and Soms (1983). In the present case (3.3) must be verified by trial and error by showing that  $\min_{\substack{\sum_{i=1}^k x_i = x_1}} \prod_{i=1}^k (n_i - x_i)^{\gamma_i} = (n_1 - x_1)^{\gamma_1} \prod_{i=2}^k n_i^{\gamma_i}$  and that

$$\max_{\substack{\sum_{i=1}^k x_i = x_1 + 1}} \prod_{i=1}^k (n_i - x_i)^{\gamma_i} < (n_1 - x_1)^{\gamma_1} \prod_{i=2}^k n_i^{\gamma_i}.$$

**Example 3.1.** Let  $k = 3$ ,  $\tilde{n} = (5, 5, 5)$ ,  $\tilde{\gamma} = (3, 3, 2)$ ,  $\alpha = .10$  and  $\tilde{x} = (1, 0, 0)$ . Then

$$\min_{\substack{\sum_{i=1}^3 x_i = 1}} \prod_{i=1}^3 (n_i - x_i)^{\gamma_i} = 200000 \quad \text{and} \quad \max_{\substack{\sum_{i=1}^3 x_i = 2}} \prod_{i=1}^3 (n_i - x_i)^{\gamma_i} = 140625 \quad \text{and}$$

$\tilde{x}$  is a key test result and (3.3) is satisfied and hence

$$a_{g(\tilde{x})} = .4161^3 = .0720,$$

where  $.10 = I_{.4161}(4, 2)$ . Further, it can also be verified that  $\tilde{x} = (2, 0, 0)$  is a key test result for which (3.3) is satisfied, but that for  $\tilde{x} = (3, 0, 0)$ , (3.3) is violated.

Note that Theorem 3.2 asserts that  $a_{g(\tilde{x})} = b_{\alpha}^{\gamma_1}$  for  $0 < \alpha < 1$ . It is thus possible that (3.3) is not true but the conclusion still holds for  $\alpha$  of practical importance. This is taken up in Section 4.

#### 4. THE LINDSTROM-MADDEN METHOD FOR SERIES SYSTEMS WITH

##### REPEATED COMPONENTS

When  $\gamma_1 = \gamma_2 = \dots = \gamma_r = 1$ , the Lindstrom-Madden method (henceforth abbreviated L-M) is an approximation  $b_{g(\tilde{x})}$  to  $a_{g(\tilde{x})}$  of the form

$$b_{g(\tilde{x})} = \min_{1 \leq i \leq k} b_{\alpha}(n_i), \quad (4.1)$$

where

$$\alpha = I_{b_{\alpha}(n_i)}(n_i - t_{0i}, t_{0i} + 1) , \quad (4.2)$$

with

$$t_{0i} = n_i \left( 1 - \prod_{j=1}^k (n_j - x_j) / n_i \right) , \quad (4.3)$$

i.e.,  $t_{0i}$  is the maximum of the recursive indexes  $t_i$  defined by (2.4). For the usual levels of  $\alpha$ ,  $b_{g(\tilde{x})} = b_{\alpha}(n_i)$ . Further, numerical evidence indicates (Harris and Soms (1983)) that for  $\alpha$  levels of practical significance

$$b_{g(\tilde{x})} < a_{g(\tilde{x})} . \quad (4.4)$$

(4.4) was incorrectly claimed to be true for  $0 < \alpha < 1$  in Sudakov (1974) and this is discussed at length in Harris and Soms (1983). However, (4.4) is known to hold for special cases (Winterbottom (1974) and Harris and Soms (1983)).

Motivated by the above, we now give an L-M approximation  $b_{g(\tilde{x})}$  to  $a_{g(\tilde{x})}$  for arbitrary  $\gamma_i$  by

$$b_{g(\tilde{x})} = \min_{1 \leq i \leq k} b_{\alpha}(n_i)^{\gamma_i} , \quad (4.5)$$

where

$$\alpha = I_{b_{\alpha}(n_i)}(n_i - t_{0i}, t_{0i} + 1) , \quad (4.6)$$

with

$$t_{0i} = n_i - \left( \prod_{j=1}^k (n_j - x_j)^{\gamma_j} \right) / \left( \prod_{\substack{j=1 \\ j \neq i}}^k n_j^{\gamma_j} \right)^{1/\gamma_i} , \quad (4.7)$$

i.e.,  $t_{0i}$  is the maximum of the recursive indexes  $t_i$  defined by (2.4). However, in this case it is not clear which index  $i$  gives the minimum, except that the likely

candidate is the one for which  $y_j$ ,  $1 \leq j \leq k$ , is a maximum. We might expect, by analogy, that for  $\alpha$  levels of practical interest

$$b_g(\tilde{x}) \leq a_g(\tilde{x}) \quad (4.8)$$

### 5. NUMERICAL EXAMPLES

For  $k = 2$  and selected  $\tilde{n}$ ,  $\tilde{y}$ ,  $\tilde{x}$ ,  $\alpha = .05$  and  $.10$ , Table I gives  $b_g(\tilde{x})$ ,  $a_g(\tilde{x})$  and the best upper bound,  $u_g(\tilde{x})$ ,

$$u_g(\tilde{x}) = \min_{1 \leq i \leq k} u_{\alpha}(n_i)^{y_i}, \quad (5.1)$$

where

$$\alpha = I_{u_{\alpha}(n_i)}(n_i - [t_{0i}], [t_{0i}] + 1) \quad (5.2)$$

and  $t_{0i}$  are defined as in (4.6).

TABLE I.

L-M Approximations and  $a_g(\tilde{x})$

$(n_1, n_2)$	$(Y_1, Y_2)$	$(x_1, x_2)$	$\alpha$	$b_g(\tilde{x})$	$a_g(\tilde{x})$	$u_g(\tilde{x})$
(10, 10)	(1, 2)	(0, 1)	.05	.3670	.3670	.3670
(10, 10)	(1, 2)	(0, 1)	.10	.4398	.4398	.4398
(10, 10)	(1, 2)	(1, 1)	.05	.3045	.3514	.3670
(10, 10)	(1, 2)	(1, 1)	.10	.3715	.4227	.4398
(10, 10)	(1, 2)	(2, 1)	.05	.2484	.3151	.3670
(10, 10)	(1, 2)	(2, 1)	.10	.3088	.3825	.4398
(10, 15)	(2, 3)	(0, 1)	.05	.3695	.3719	.3742
(10, 15)	(2, 3)	(0, 1)	.10	.4425	.4446	.4467
(10, 15)	(2, 3)	(1, 1)	.05	.2554	.3042	.3670
(10, 15)	(2, 3)	(1, 1)	.10	.3167	.3705	.4398
(10, 15)	(2, 3)	(2, 1)	.05	.1712	.1981	.2431
(10, 15)	(2, 3)	(2, 0)	.10	.2203	.2513	.3029

Note that for all the cases in Table I,  $b_{\tilde{q}}(\tilde{x})$  is a lower bound for  $a_{\tilde{q}}(\tilde{x})$ . The computations were done by a short FORTRAN program, a listing of which can be obtained from the author.

#### 6. CONCLUDING REMARKS

In this paper we have extended the L-M method to series systems with repeated components. More work is needed to ascertain the region of validity of (4.8).

# BIBLIOGRAPHY

Buehler, R. J., (1957). Confidence Limits for the Product of Two Binomial Parameters.

Journal of the American Statistical Association, 52, 482-93.

Harris, B. and Soms, A. P., (1983). The Theory of Optimal Confidence Limits for Systems Reliability with Counterexamples for Results on Optimal Confidence Limits for Series Systems. Technical Report #708, Department of Statistics, University of Wisconsin-Madison.

Pledger, G. and Proschan, F., (1971). Comparison of Order Statistics and of Spacings from Heterogeneous Distributions. Optimizing Methods in Statistics, New York: Academic Press, 89-113.

Sudakov, R. S., (1974). On the Question of Interval Estimation of the Index of Reliability of a Sequential System. Engineering Cybernetics, 12, 55-63.

Winterbottom, A., (1974). Lower Limits for Series System Reliability from Binomial Data. Journal of the American Statistical Association, 69, 782-8.

APS:scr

REPORT DOCUMENTATION PAGE		READ INSTRUCTIONS BEFORE COMPLETING FORM
1. REPORT NUMBER  2732	2. GOVT ACCESSION NO.	3. RECIPIENT'S CATALOG NUMBER
4. TITLE (and Subtitle)  THE LINDSTROM-MADDEN METHOD FOR SERIES SYSTEMS WITH REPEATED COMPONENTS		5. TYPE OF REPORT & PERIOD COVERED Summary Report - no specific reporting period
		6. PERFORMING ORG. REPORT NUMBER
7. AUTHOR(s)  Andrew P. Soms		8. CONTRACT OR GRANT NUMBER(s)  DAAG29-80-C-0041 N00014-79-C-0321
9. PERFORMING ORGANIZATION NAME AND ADDRESS Mathematics Research Center, University of 610 Walnut Street Madison, Wisconsin 53706		10. PROGRAM ELEMENT, PROJECT, TASK AREA & WORK UNIT NUMBERS Work Unit Number 4 - Statistics and Probability
11. CONTROLLING OFFICE NAME AND ADDRESS  See Item 18 below.		12. REPORT DATE August 1984
		13. NUMBER OF PAGES 12
14. MONITORING AGENCY NAME & ADDRESS (if different from Controlling Office)		15. SECURITY CLASS. (of this report)  UNCLASSIFIED
		15a. DECLASSIFICATION/DOWNGRADING SCHEDULE
16. DISTRIBUTION STATEMENT (of this Report)  Approved for public release; distribution unlimited.		
17. DISTRIBUTION STATEMENT (of the abstract entered in Block 20, if different from Report)		
18. SUPPLEMENTARY NOTES U. S. Army Research Office P. O. Box 12211 Research Triangle Park North Carolina 27709  Office of Naval Research 800 North Quincy Street Arlington, VA 22217		
19. KEY WORDS (Continue on reverse side if necessary and identify by block number)  Lindstrom-Madden approximation; Optimal confidence limits; Reliability; Repeated components; Series system		
20. ABSTRACT (Continue on reverse side if necessary and identify by block number)  The Lindstrom-Madden method of computing lower confidence limits for series systems with unlike components is extended to series systems with repeated components utilizing the results of Harris and Soms (1983). An exact solution is given for no failures and key test results, together with an approximation for the general case. Numerical examples are also provided.		